G 364: Mobile and Wireless Networking CLASS 13, Mon. Feb. 23 2004 Stefano Basagni Spring 2004 M-W, 11:40am-1:20pm, 109 Rob

Why New Protocols for Sensor Networks?, 1

- Not just MANET with static nodes!
- Several issues
 - Severe energy constraints
 - Large scale
 - Potential to exploit mobility
 - Applications involve real-time constraints and control loops

Why New Protocols for Sensor Networks?, 2

- Purpose is to estimate function of a physical phenomenon, and not just to move bits
 - E.g. entire map, average, some event etc.
 - Don't really care *which* node provides that
 - Three implications:
 - no longer name or address nodes
 - collaborative computation instead of communication
 - data correlation can be exploited

Wireless Sensor Network Protocols

Building long-lived, massively-distributed, physically-coupled systems: Coordinating to minimize duty cycle and communication Adaptive MAC Adaptive Topology Routing In-network processing Data centric routing Programming models

MAC in Sensor Nets

- Important attributes of MAC protocols
 - 1. Collision avoidance
 - 2. Energy efficiency
 - 3. Scalability in node density
 - 4. Latency
 - 5. Fairness
 - 6. Throughput
 - 7. Bandwidth utilization

MAC Impact on Networks

- Major sources of energy waste
 - Idle listening when no sensing events, Collisions, Control overhead, Overhearing



Identifying the Energy Consumers



Energy Efficiency in MAC

- Major sources of energy waste
 - Idle listening
 - Long idle time when no sensing event happens
 - Collisions
 - Control overhead
 - Overhearing

- Common to all
- wireless networks
- Try to reduce energy consumption from all above sources
- TDMA requires slot allocation and time synchronization
- Combine benefits of TDMA + contention protocols

Sensor-MAC (S-MAC) Design (Wei et al. 2002)

Tradeoffs



- Major components of S-MAC
 - Periodic listen and sleep
 - Collision avoidance
 - Overhearing avoidance
 - Message passing

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Collision Avoidance

- Problem: Multiple senders want to talk
 Options: Contention vs. TDMA
 Solution: Similar to IEEE 802.11 ad hoc mode (DCF)
 - Physical and virtual carrier sense
 - Randomized backoff time
 - RTS/CTS for hidden terminal problem
 RTS/CTS/DATA/ACK sequence

Overhearing Avoidance

- Problem: Receive packets destined to others
- Solution: Sleep when neighbors talk
 - Basic idea from PAMAS (Singh 1998)
 - But we only use in-channel signaling
- Who should sleep?
 - All immediate neighbors of sender and receiver

How long?

The *duration* field in each packet informs other nodes the sleep interval

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Testbed Nodes

- Platform
 - Motes (UC Berkeley) :
 - 8-bit CPU at 4MHz,
 8KB flash, 512B RAM
 - TinyOS: avont drivon
 - TinyOS: event-driven
 - Compared MAC modules
 - IEEE 802.11-like protocol



- Message passing with overhearing avoidance
- S-MAC (2 + periodic listen/sleep)
- URL: http://www.isi.edu/scadds/smac/

Adaptive Topology (AT)

- Can we do more than shut down radio in between transmissions/receptions?
- Can we put nodes to sleep for longer periods of time?
- Goal:
 - Exploit high density (over) deployment to extend system lifetime
 - Provide topology that adapts to the application needs
 - Self-configuring system that adapts to environment without manual configuration

AT: Problem Description

Simple Formulation (Geometric Disk Covering)

- Given a distribution of **N** nodes in a plane
- Place a *minimum* number of disks of radius *r* (centered on the nodes) to cover them
- Disk represents the radio connectivity (*simple* circle model)
- The problem is NP-Hard

Connectivity Measurements*

There is a *non-zero* probability of receiving packets at distances much greater than the average cell range

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Can't just determine connectivity clusters thru geographic coordinates

For the same reason you can't determine coordinates w/connectivity

*An Empirical Study of Epidemic Algorithms in Large Scale Multihop Wireless Networks Ganesan, Krishnamachari, Woo, Culler, Estrin and Wicker, UCLA/CSD-TR 02-0013. 2/23/04

Tradeoff

How many nodes to activate?

- Few active nodes:
 - distance between neighboring nodes high → increased packet loss, higher transmit power and reduced spatial reuse
 - need to maintain sensing coverage
- Too many active nodes:
 - at best, expending unnecessary energy
 - at worst nodes may *interfere* with one another by *congesting* the channel

Adaptive Topology Schemes

Mechanisms being explored:

- *Empirical adaptation*: Each node assesses its connectivity and adapts participation in multi-hop topology based on the measured operating region, ASCENT (Cerpa et al. 2002)
- Cluster-based, load sharing within clusters, CEC (Xu et al. 2002)
- Routing/Geographic topology based, eliminate redundant links, SPAN (Chen et al. 2001), GAF (Xu et al. 2001)
- Data/traffic driven: Trigger nodes on demand using paging channel, STEM (Tsiatsis et al. 2002)

Topology Control

Deciding on:

- which nodes turn on
- when they turn on, and
- at what Tx power

So that **desired network connectivity** is maintained

Motivation for Topology Control, 1



- High power
- High interference
- Low Throughput

Motivation for Topology Control, 2



- Low power
- Low interference
- High Throughput
- Global Connectivity

GAF: Geographic Adaptive Fidelity

Geography-informed Energy
 Conservation for Ad Hoc Routing

 Ya Xu, John Heidemann, Deborah Estrin USC/ISI, UCLA

Energy \Leftrightarrow Node Density



GAF: Geographic Adaptive **Fidelity**



GAF reduces the energy by a factor M' This factor is a function of the average number of nodes in a grid: M

 $M' = \frac{M}{1 - e^{-M}} \qquad for uniformly random$ node deployment

Energy \Leftrightarrow Density

Conserve traffic forwarding capacity Divide network in virtual grids of size r <= radio range/sqrt(5) Each node in a grid is equivalent from a traffic forwarding perspective

Keep 1 node awake in each grid at each time

			n
M'	М	λ	eigl
1.0	0	0	nboi
1.5	0.87	13.7	nur s of
2.0	1.59	25.0	nbei a n
2.5	2.22	35.0	r of ode
3.0	2.82	44.3	23

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Topology Control: Other solutions

- ASCENT: Adaptive Self-Configuring sEnsor Networks Topologies
 - Cerpa & Estrin, UCLA
- Span: An Energy-Efficient Coordination Algorithm for Topology Maintenance in Ad Hoc Networks
 - Chen, Jamieson, Balakrishnan & Morris, MIT
- STEM: Sparse Topology and Energy Management
 - Schurgers, Tsiatsis, Ganeriwal & Srivastava, UCLA

Routing in Sensor Networks

Given a topology, how to route data?

- The MANET way: Reactive, proactive, and geo-enabled routing ...
- Building on Geo Routing
 GRAB (Lu et al 2002)

Routing on curve (Badri 2002)

Directed Diffusion (DD): Data Centric Routing, 1

Basic idea

- name data (not nodes) with externally relevant attributes
 - Data type, time, location of node, SNR, etc
- diffuse requests and responses across network using application driven routing (e.g., geo sensitive or not)
- support in-network aggregation and processing
- optimize path with gradient-based feedback

DD, 2

Data sources publish data, data clients subscribe to data

- However, all nodes may play both roles
 - A node that aggregates/combines/processes incoming sensor node data becomes a source of new data
 - A sensor node that only publishes when a combination of conditions arise, is a client for the triggering event data
- True peer to peer system

Linux (32 bit proc) and TinyOS (8 bit proc) implementations

Assignments

Download the survey on sensor nets

Updated information on the class web page:

www.ece.neu.edu/courses/eceg364/2004sp

